

PRIVACY-PRESERVING SIGNATURES FROM ISOGENIES

Federico Pintore

Department of Mathematics, University of Bari, IT

Joint work with Ward Beullens¹ and Schuichi Katsumata²

¹imec-COSIC, BE

²National Institute of Advanced Industrial Science and Technology (AIST), JP

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RING SIGNATURES

Introduced in 2001 by Rivest, Shamir, Tauman.

Proposed application: **whistleblowers**



Privacy-preserving cryptocurrencies: Dash, **Monero**, ZCash

The security of the ring signature used by Monero relies on the hardness of the DLP.

CALAMARI AND FALAFL

- A simple and efficient **generic construction** of OR-proofs from simple algebraic structure (i.e., "admissible" group actions).
- The OR-proof is converted into a **ring signature**, where the signature size scales **logarithmically in N** , better than previously known PQ schemes ($c \cdot 2\lambda \cdot \log_2 N$).

Calamari

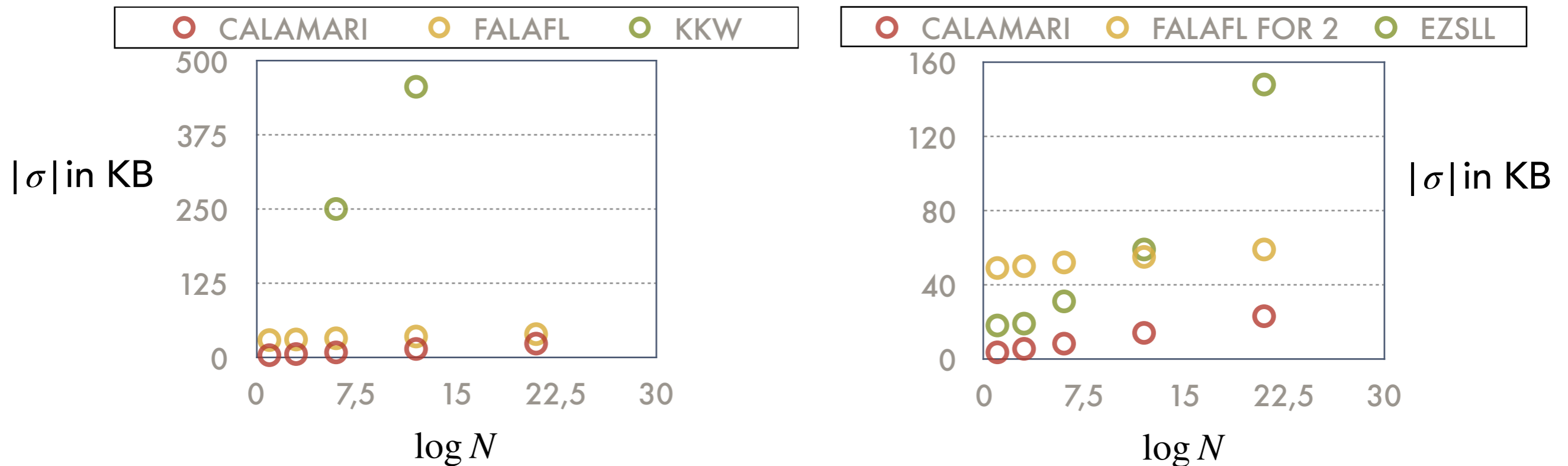
- **first** isogeny-based ring signature
- **smallest signatures** among PQ alternatives

Falafel

- **fast signing**
- smaller signature size than lattice-based alternatives from $N \approx 1024$

POST-QUANTUM RING SIGNATURES

Ring signatures having size that scales logarithmically in N have been proposed only from **symmetric-key** primitives, and **lattice-based** assumptions.



- **FalafI:** $|\sigma| = 0.5 \log N + 29$ KB
- **Raptor:** $|\sigma| = 1.3 N + 1.5$ KB

ROADMAP

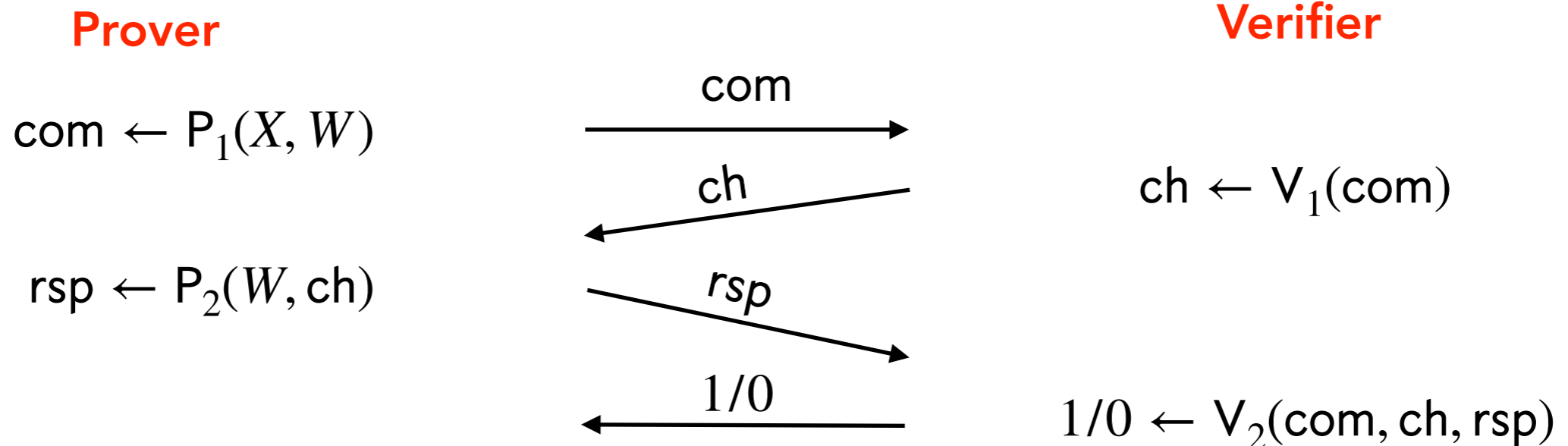
1. OR-proofs
2. Admissible Group Actions
3. A new OR-Proof
4. Calamari and Falafel

SIGMA PROTOCOLS

Let $R \subset X \times W$ be a binary relation. A sigma protocol

$$\Pi_{\Sigma} = (P = (P_1, P_2), V = (V_1, V_2))$$

for R is a **three-move interactive protocol** between a **prover** (holding a verification-secret key pair $(X, W) \in R$) and a **verifier** (holding X).



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Required properties

- Correctness with abort
- **(Non-abort) Honest-Verifier Zero-Knowledge**
- Special Soundness
- ...

OR-PROOFS & RING SIGNATURES


An **OR-Proof** for the binary relation $R \subset X \times W$ is a sigma protocol for

$$R_{\text{OR}} = \{((X_1, \dots, X_N), (W, I)) \mid N \in \mathbb{N}, I \in [N], (X_I, W) \in R\}$$

An OR-Proof, with negligibly-small soundness error, can be turned into a **ring signatures** by means of the Fiat-Shamir transform.

- $\text{ch} = H(\text{com}, R, m)$
- $\sigma = (\text{com}, \text{ch}, \text{rsp})$

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ADMISSIBLE GROUP ACTIONS (AGAs)

A **AGA** with respect to $X_0 \in X$ is a tuple $(\mathbb{G}, X, S_1, S_2, \delta, D_X)$, where:

- \mathbb{G} is a **group** and S_1, S_2 are **symmetric subsets** of \mathbb{G} ($s \in S_i \Rightarrow -s \in S_i$);
- X is a **finite set**;
- $\delta \in [0,1]$;
- D_X is a distribution over a set of **group actions**

$$\star : \mathbb{G} \times X \rightarrow X.$$

$$(g, X) \mapsto g \star X$$

- $0_G \star X = X$;
- $g_1 \star (g_2 \star X) = (g_1 + g_2) \star X$

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► For the set $S_3 = \bigcap_{g \in S_1} \{S_2 + g\}$, it holds $|S_3| = \delta \cdot |S_2|$.

► Given $X = g \star X_0$, with g sampled uniformly in S_1 , it is **hard to find** g .

Example (DHKE on a group \mathbb{H}): $X = \mathbb{H}^*$, $\mathbb{G} = S_1 = S_2 = \mathbb{Z}_q^*$, $\delta = 1$, $\star = \exp$

NOTATION



Given an **AGA** $(\mathbb{G}, X, S_1, S_2, \delta, D_X)$ with respect to $X_0 \in X$ we denote:

- $|g|$ = bits to represent $g \in \mathbb{G}$;
- x = bits to represent $X \in X$.

The size of a transcript (com, ch, rsp) is denote by $|\text{trans}|$.

The size of a ring signature is denote by $|\sigma|$.

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CONSTRUCTING AN OR-PROOF FOR AGAs - $N=1$



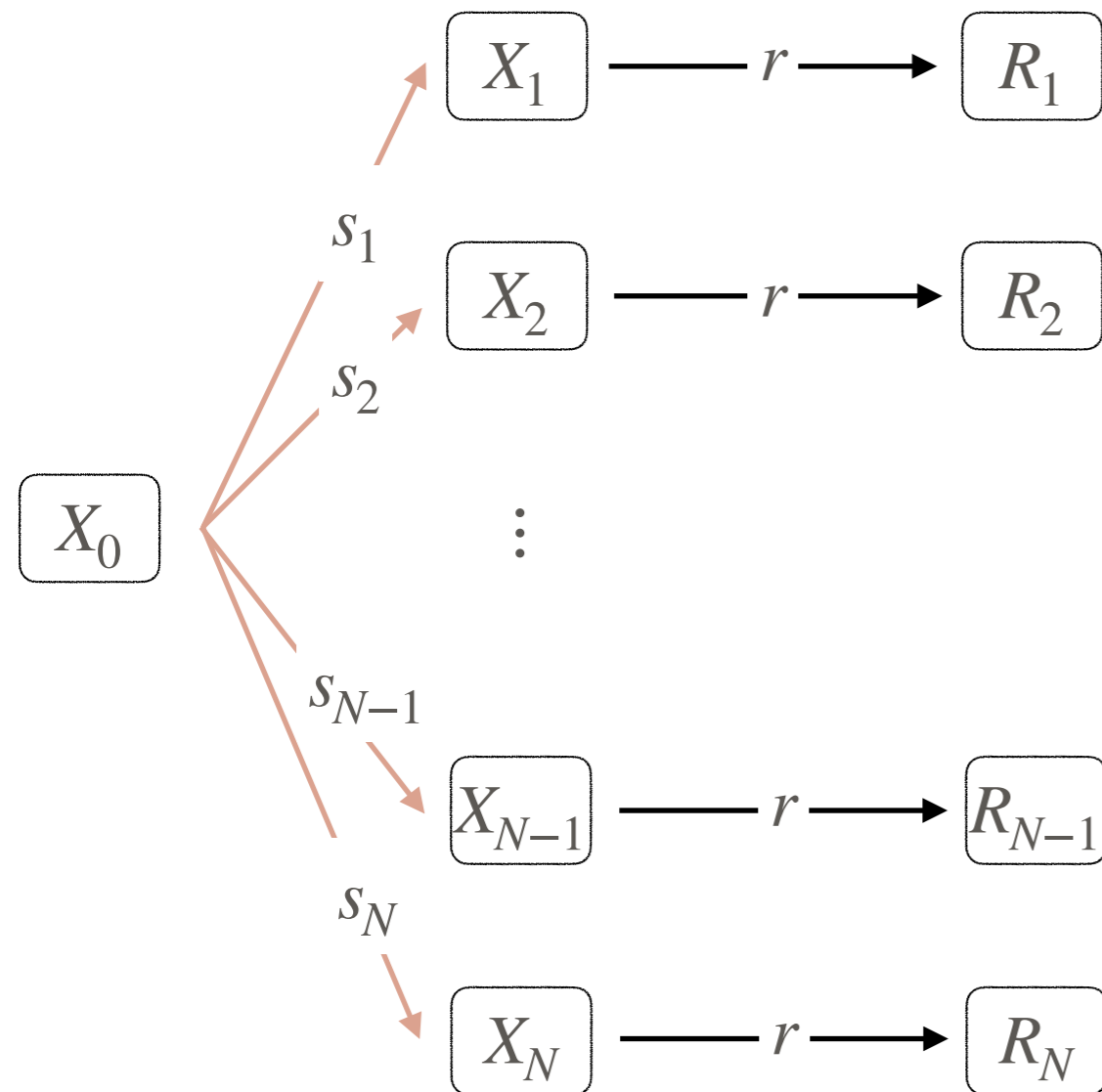
$$\text{com} = R$$

$$\text{rsp}_0 = r + s$$

$$\text{rsp}_1 = r$$

$$R = \{(X, s) \mid s \in S_1, X \in X, X = s \star X_0\}$$

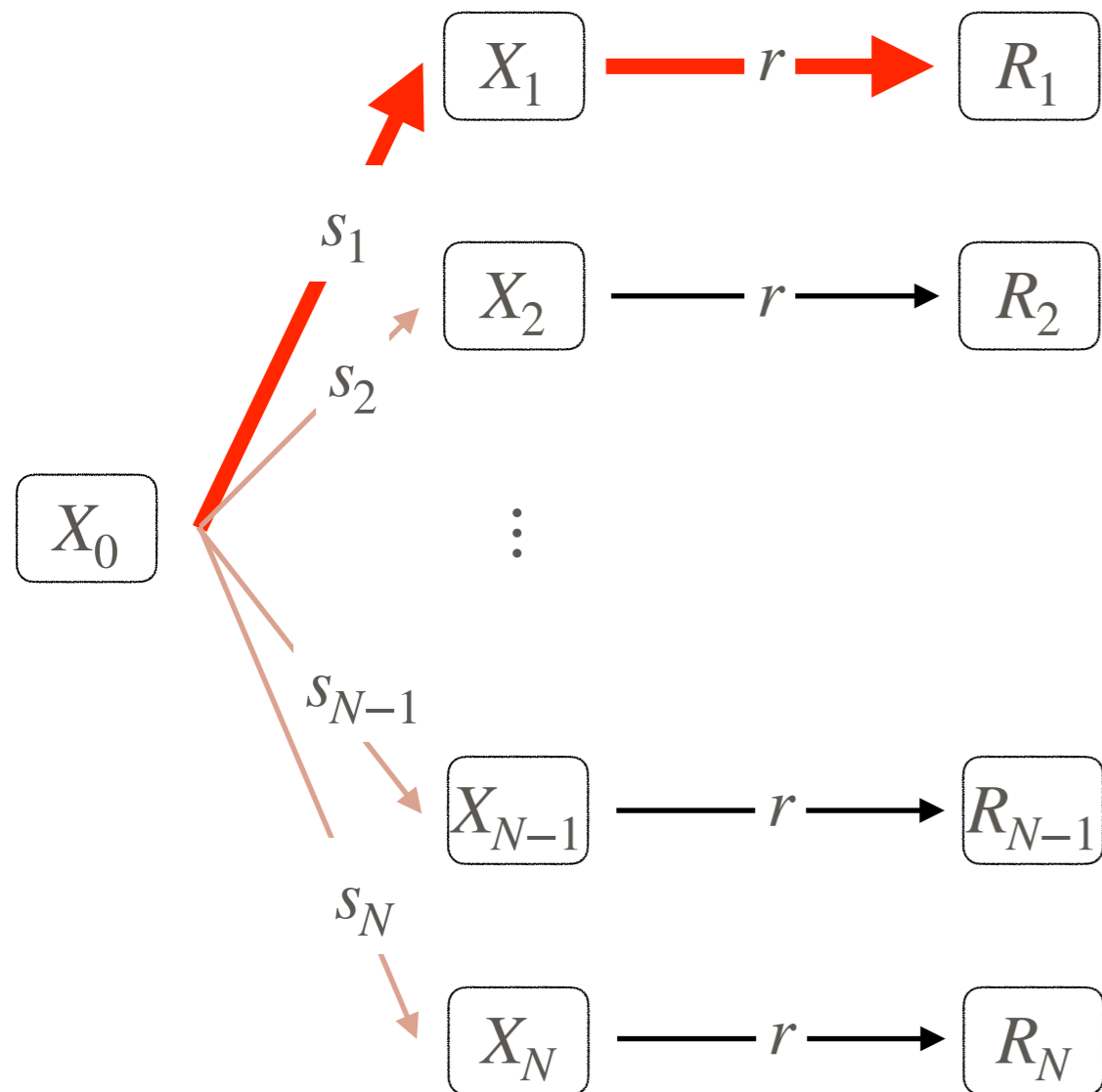
CONSTRUCTING AN OR-PROOF FOR AGAs - N



$$\text{com} = (R_1, \dots, R_N)$$

$$R_{\text{OR}} = \{((X_1, \dots, X_N), (s, I) \mid s \in S_1, X_i \in X, I \in [N], X_I = s \star X_0)\}$$

CONSTRUCTING AN OR-PROOF FOR AGAs - N




**Verification
reveals I**

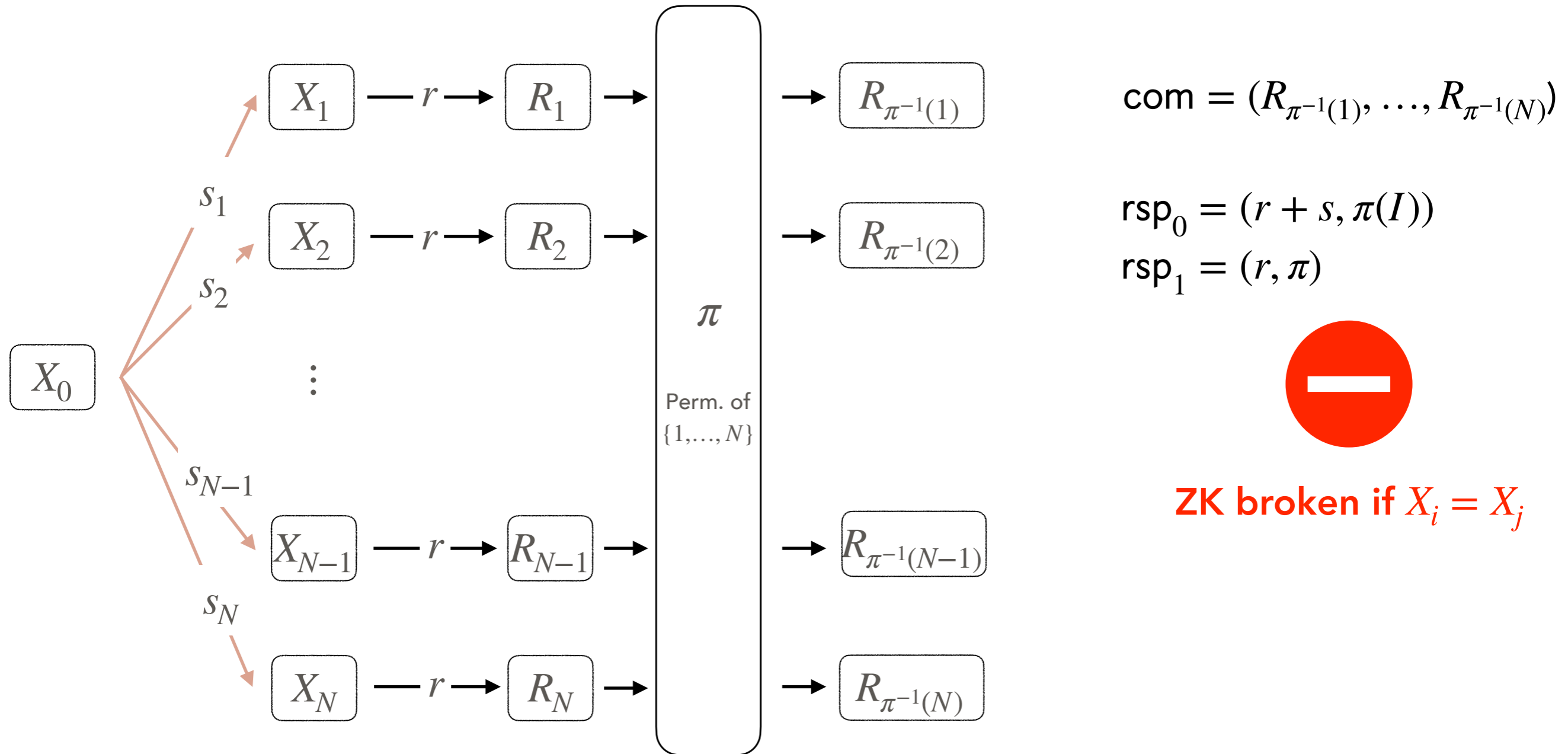
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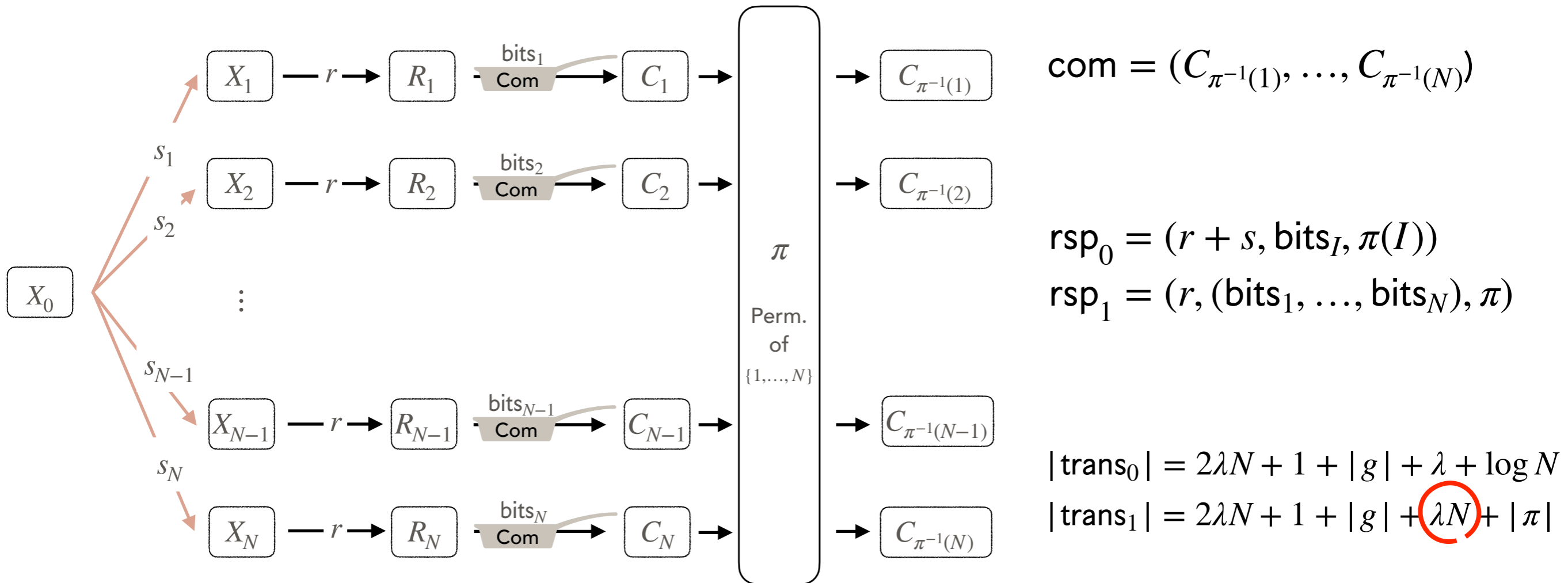
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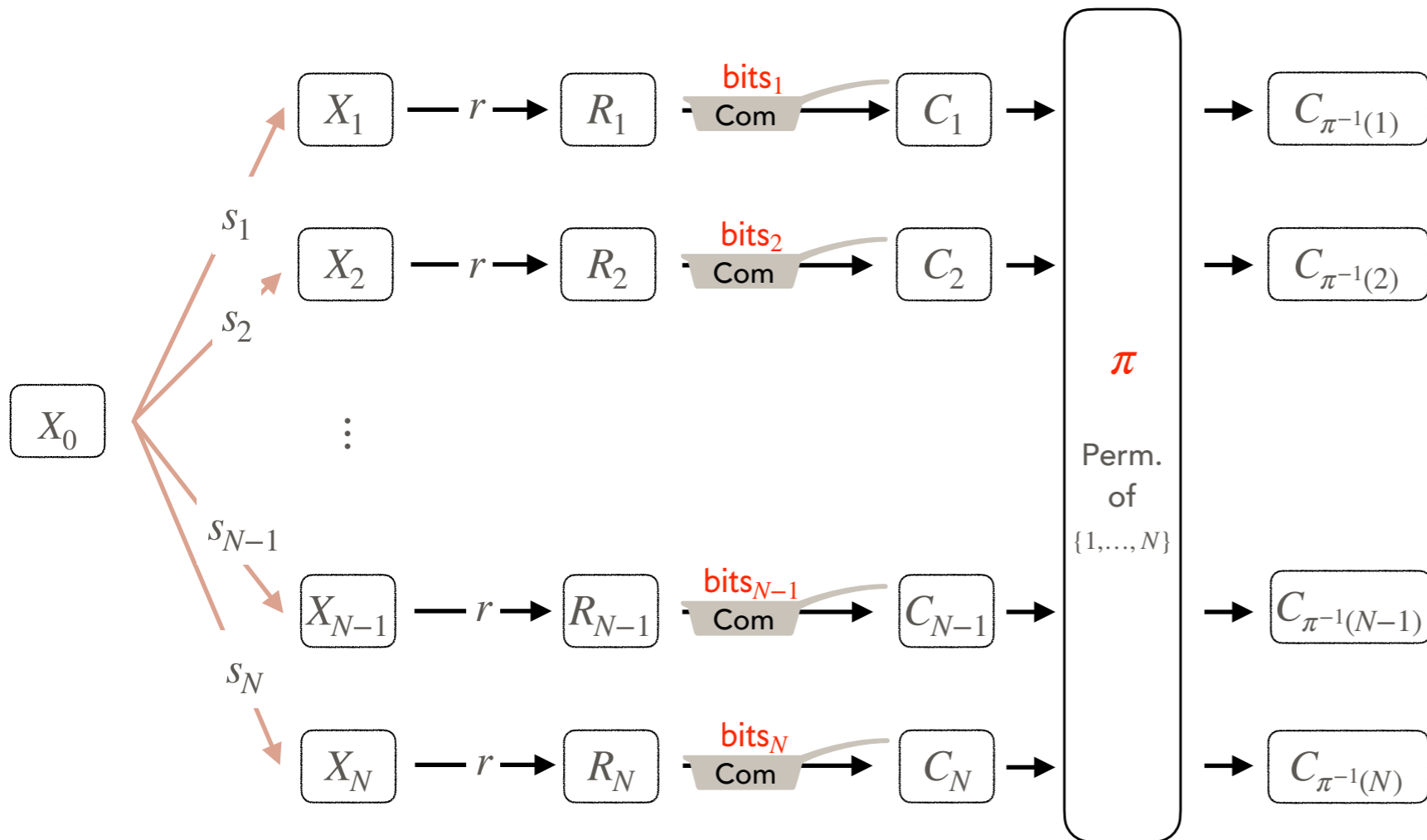
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CONSTRUCTING AN OR-PROOF FOR AGAs - N

$$(r, \pi, \text{bits}_1, \dots, \text{bits}_N) \leftarrow \text{PRG}(\text{seed})$$



$$\text{com} = (C_{\pi^{-1}(1)}, \dots, C_{\pi^{-1}(N)})$$

$$\text{rsp}_0 = (r + s, \text{bits}_I, \pi(I))$$

$$\text{rsp}_1 = (r, (\text{bits}_1, \dots, \text{bits}_N), \pi)$$

seed

$$|\text{trans}_0| = 2\lambda N + 1 + |g| + \lambda + \log N$$

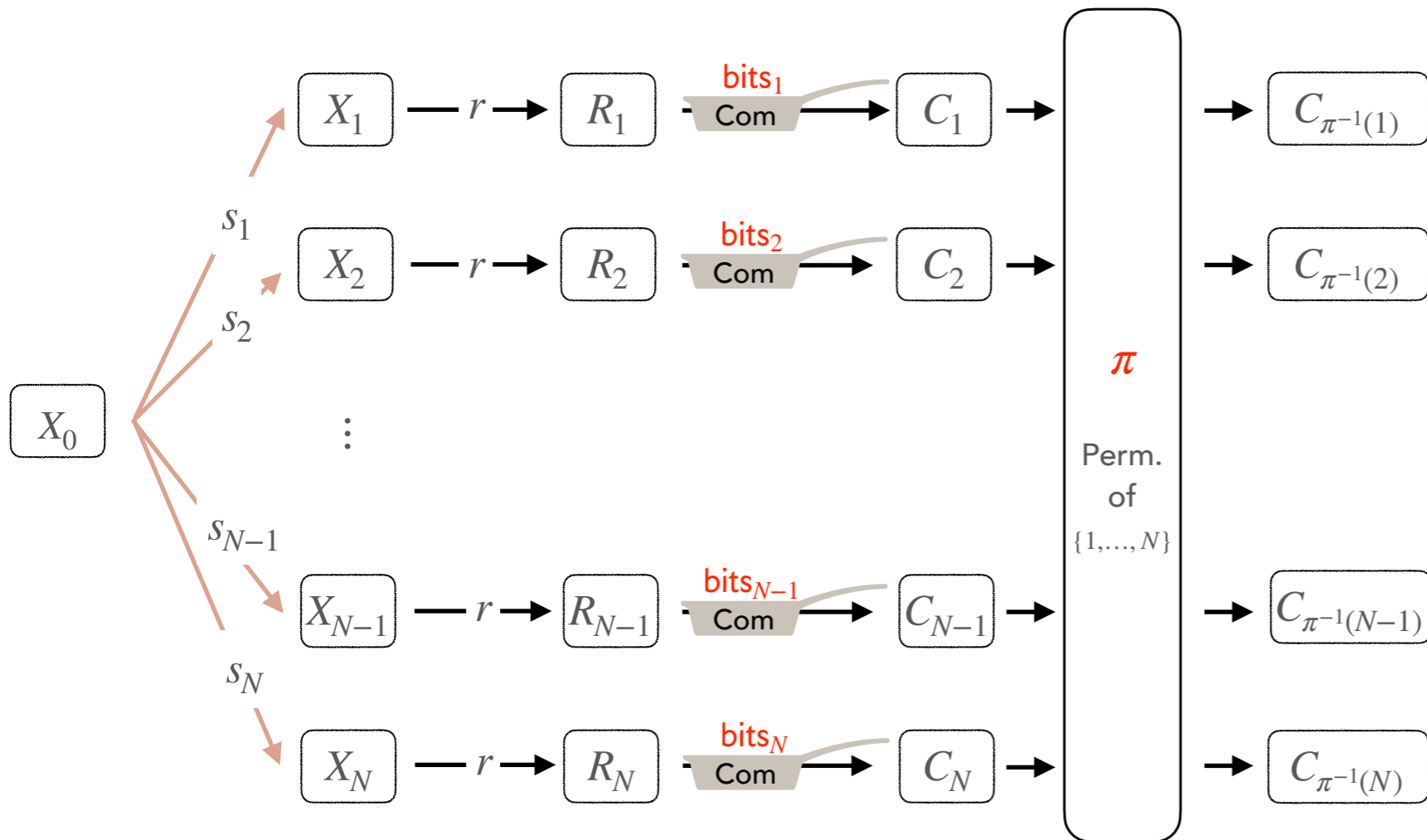
$$|\text{trans}_1| = 2\lambda N + 1 + |g| + \lambda N + |\pi|$$

λ

$$R_{\text{OR}} = \{((X_1, \dots, X_N), (s, I) \mid s \in S_1, X_i \in X, I \in [N], X_I = s \star X_0)\}$$

CONSTRUCTING AN OR-PROOF FOR AGAs - N

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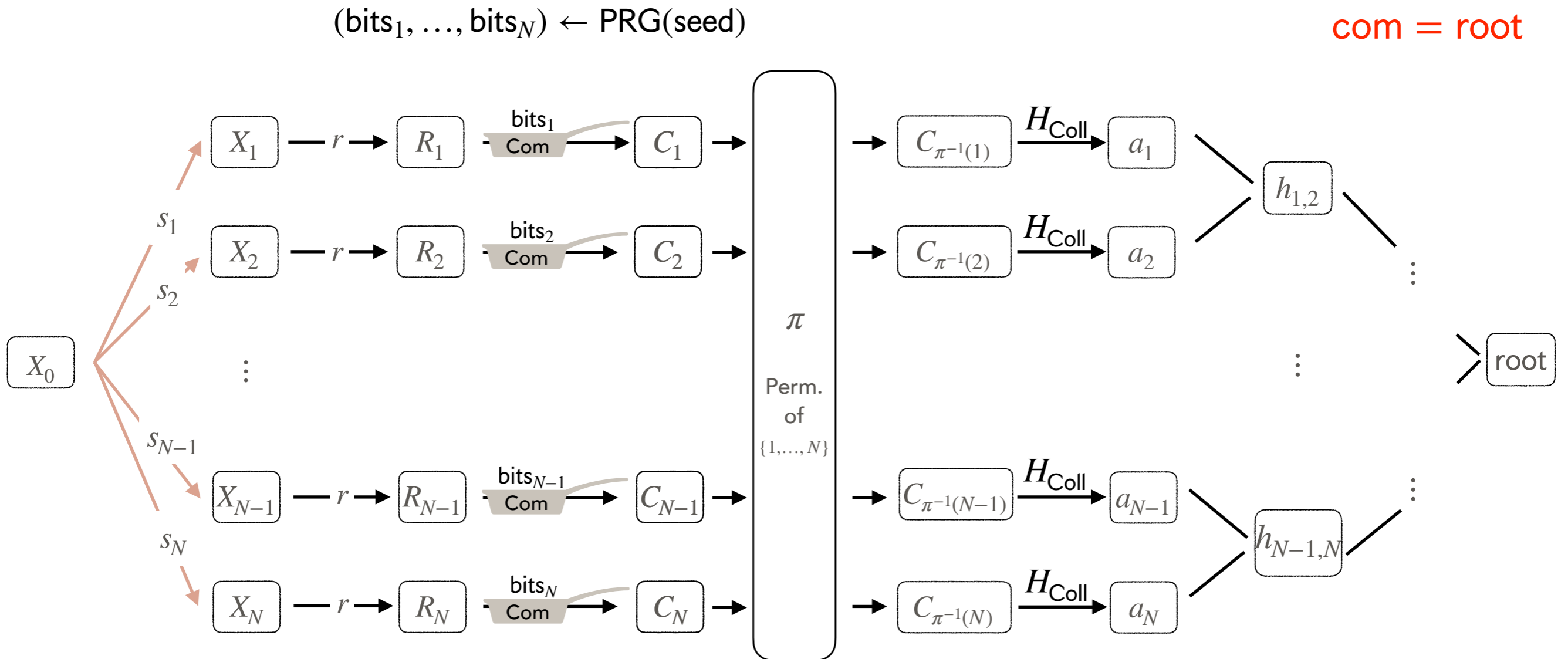
$$|\text{trans}_0| = 2\lambda N + 1 + |g| + \lambda + \log N$$

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λ

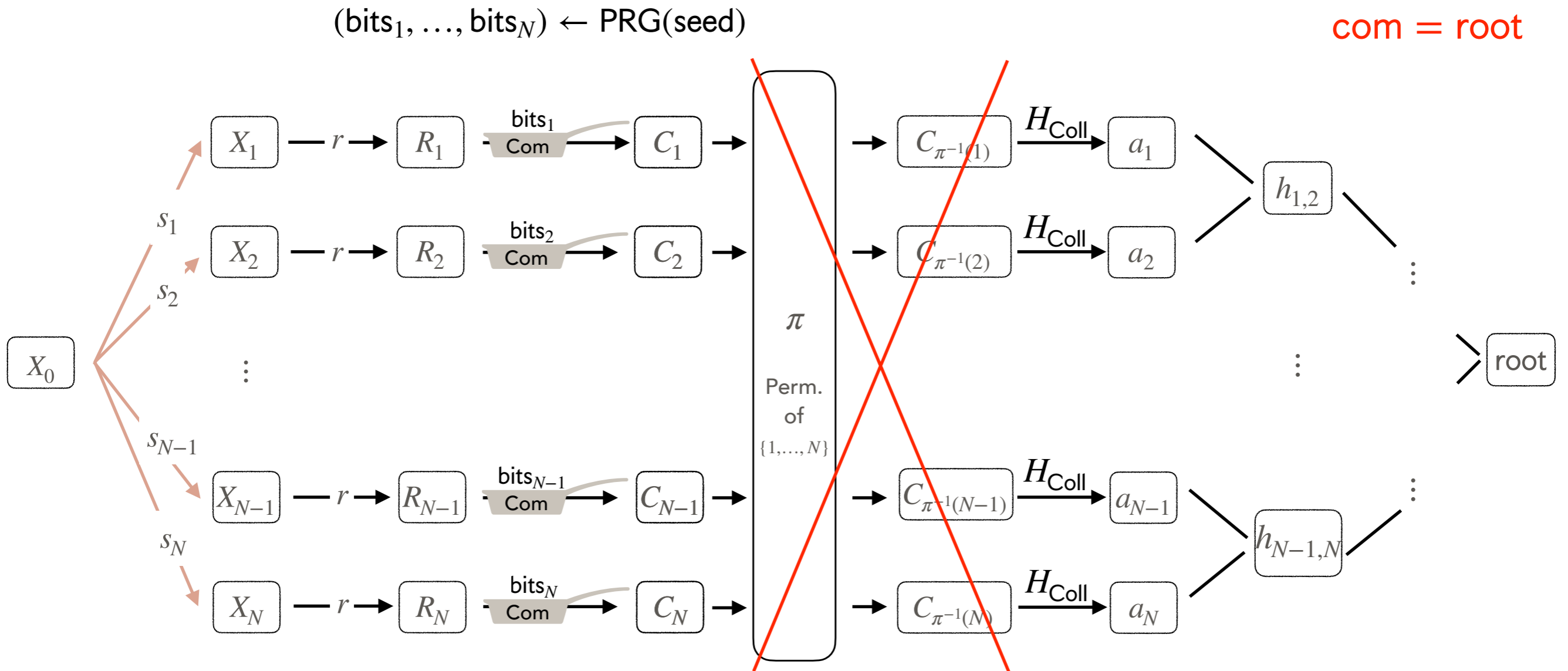
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OR-PROOF FOR AGAs - OPTIMISATION 1



$$R = \{((X_1, \dots, X_N), (s, I) \mid s \in S_1, X_i \in X, I \in [N], X_I = s \star X_0)\}$$

OR-PROOF FOR AGAs - OPTIMISATION 1

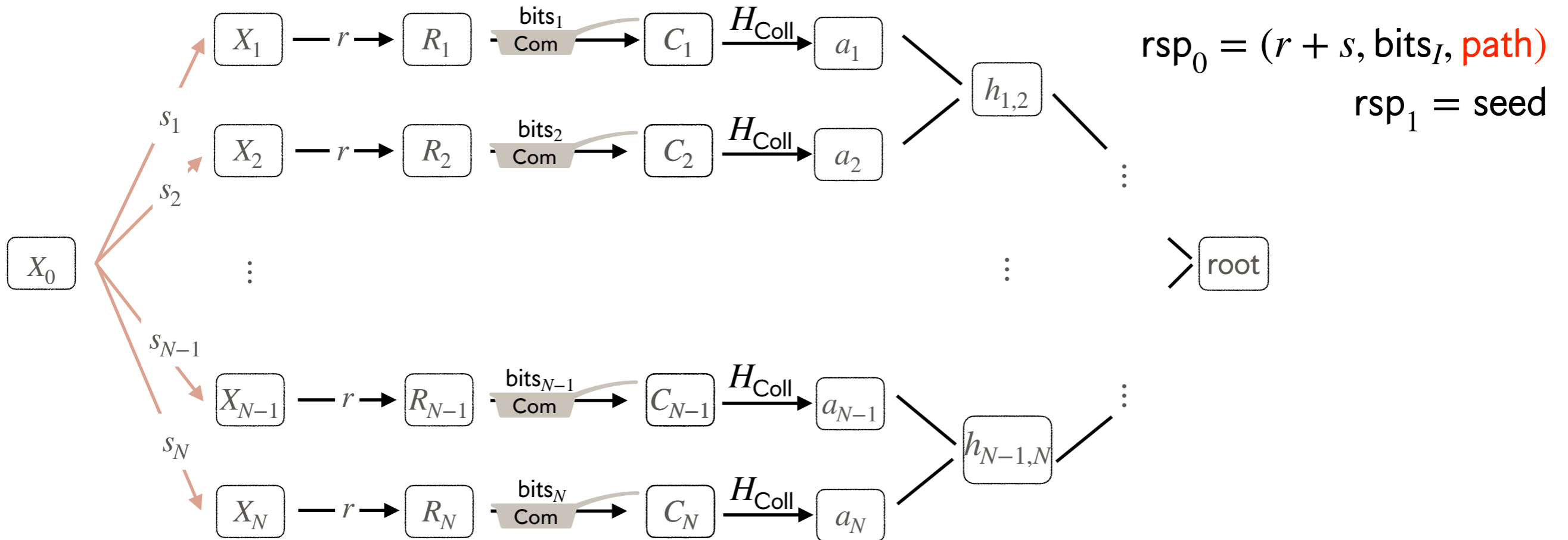


The permutation is no longer needed!
No indices involved in the verification

OR-PROOF FOR AGAs - OPTIMISATION 1

$(\text{bits}_1, \dots, \text{bits}_N) \leftarrow \text{PRG}(\text{seed})$

com = root



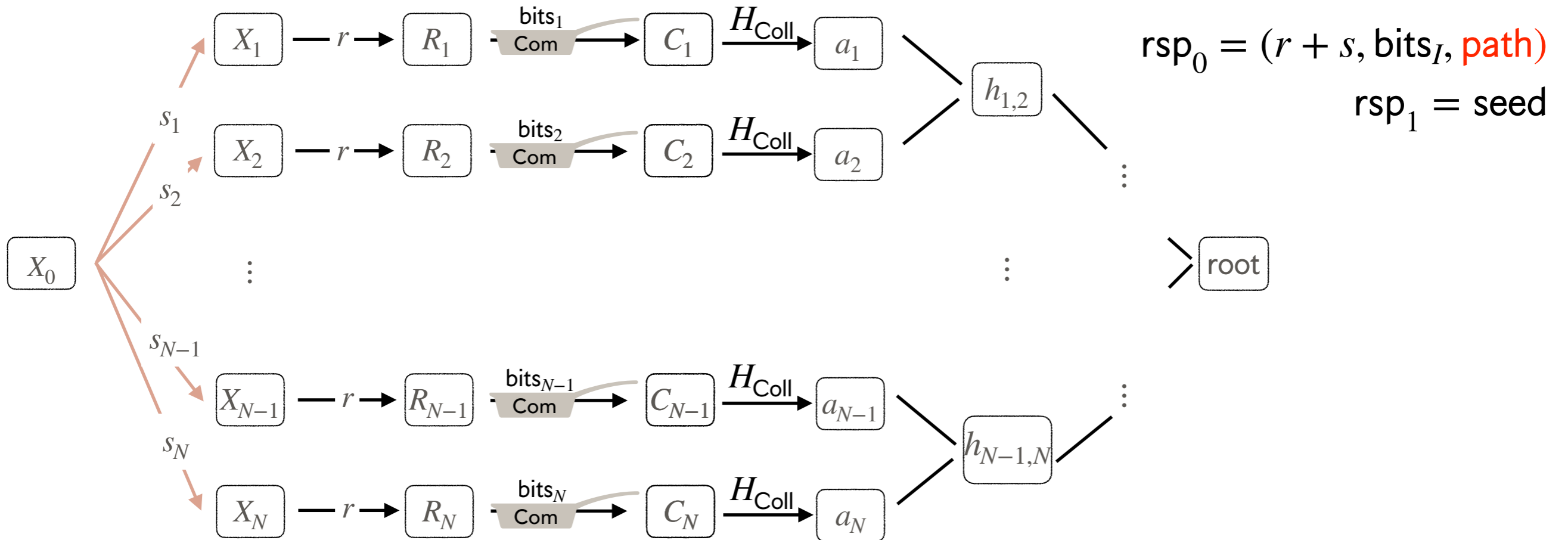
$$|\text{trans}_0| = 2\lambda + 1 + |g| + \lambda + 2\lambda(\log N - 1)$$

$$|\text{trans}_1| = 2\lambda + 1 + \lambda$$

OR-PROOF FOR AGAs - OPTIMISATION 1

$(\text{bits}_1, \dots, \text{bits}_N) \leftarrow \text{PRG}(\text{seed})$

com = root



$$|\text{trans}_0| = 1 + |g| + \lambda + 2\lambda \log N$$

$$|\text{trans}_1| = 2\lambda + 1 + \lambda$$

OR-PROOF FOR AGAs - FURTHER OPTIMISATIONS

To make the challenge error negligible small, λ parallel runs are executed.

We choose M, K such that $\binom{M}{K} \geq 2^\lambda$ and we set the challenge space to $C_{M,K} = \{\text{ch} \in \{0,1\}^M : w(\text{ch}) = M - K\}$

It could be reduced to λ




$$|\sigma| = 2\lambda + M + (M - K)\lambda + K(|g| + \lambda + 2\lambda(\log N - 1))$$

$|\text{com}|$
 $|\text{ch}|$

$|\text{rsp}_1|$

$|\text{rsp}_0|$

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CA(LA)MARI

Given a prime p , the **ideal class group** $\mathbb{G} = \text{Cl}(\mathcal{O})$ of $\mathcal{O} = \mathbb{Z} \left[\sqrt{-p} \right]$ **acts** freely and transitively on the **set X of supersingular elliptic curves** E over \mathbb{F}_p s.t. $\text{End}_p(E) \simeq \mathcal{O}$.

GAIP problem: given $X_0 \in X$ and a uniform $X \in X$, find $a \in \mathbb{G}$ s.t. $a \star X_0 = X$.

Only for the set of parameters CSIDH-512 the structure of \mathbb{G} is known.

- $\lambda = 128$
- $M = 247, K = 30$
- $|g| = 258$ ($|\mathbb{G}| \approx 2^{257.1}$)

Experimental results ($N = 8$): $|\sigma| \approx \log N + 3.5\text{KB}$

Signing ≈ 79 s

Thanks for your attention

Federico Pintore

Department of Mathematics, University of Bari (IT)

federico.pintore@gmail.com